# Victoria University of Wellington <br> School of Mathematics and Statistics 

MATH $361 \quad$ Assignment 1 T1 2024

Due 3pm Monday 11 March

1. Prove Exercise 1.2: "Let $G$ be a graph with a walk $W$ from a vertex $u$ to a vertex $v$. Then there is a path from $u$ to $v$ that uses a subset of the edges of $W$."
2. The complement of a simple graph $G=(V, E)$ is the simple graph $\bar{G}=(V, \bar{E})$, where an edge $x y$ (for any distinct $x, y \in V$ ) is in $\bar{E}$ if and only if $x y$ is not in $E$. A simple graph is self-complementary if it is isomorphic to its complement. Prove:
(a) If $G$ is a disconnected graph, then $\bar{G}$ is connected.
(b) Every non-empty self-complementary graph is connected. (Hint: use (a).)
(c) If $G$ is self-complementary, then either $|V| \equiv 0 \bmod 4$ or $|V| \equiv 1 \bmod 4$.
3. Let $s$ and $t$ be positive integers with $s \leq t$. Recall that $P_{s}$ is the path graph on $s$ vertices. Give a formula for the minimum number of edges that need to be removed from $K_{t}$ so that it has a graph isomorphic to $P_{s}$ as an induced subgraph.
4. The distance $d(u, v)$ between two vertices in a graph $G$ is defined as the length of the shortest path that joins $u$ and $v$. Prove that the distance satisfies the triangle inequality, that is, prove that $d(u, w) \leq d(u, v)+d(v, w)$ for any three vertices $u, v$ and $w$ of $G$.
5. Let $u$ be a vertex of odd degree in the graph $G$. Prove that there is a path from $u$ to another vertex of odd degree in $G$.
6. We know that trees with at least two vertices have at least two leaves. But typically trees have more leaves than that.
(a) Show that if a tree has a vertex of degree $k$, then it has at least $k$ leaves.
(b) Let $T$ be a tree with $n$ vertices, $k$ leaves, and a vertex with degree $k$, where $k \geq 2$. Suppose that $n>k+1$. Prove that $T$ has a vertex of degree two.
7. A graph is $k$-regular if every vertex has degree $k$. Prove or disprove the following:
(a) If $G$ is a $k$-regular bipartite graph, with $k \geq 2$, then $G$ has no bridges.
(b) If $G$ is a $k$-regular graph, with $k \geq 2$, then $G$ has no bridges.

## Tutorial exercises:

1. Prove Theorem 1.5: "For any graph $G$ with vertex set $V$, the relation $\sim$ is an equivalence relation on $V$."
2. Prove Theorem 1.6: "For each positive integer $n$, the complete graph $K_{n}$ has $n(n-1) / 2$ edges."
3. The complement of a simple graph $G=(V, E)$ is the simple graph $\bar{G}=(V, \bar{E})$, where an edge $x y$ (for any distinct $x, y \in V$ ) is in $\bar{E}$ if and only if $x y$ is not in $E$. A simple graph is self-complementary if it is isomorphic to its complement.
(a) Give an example of a graph that is self-complementary.
(b) Prove that every self-complementary graph with $4 k+1$ vertices has a vertex of degree $2 k$ (where $k$ is a non-negative integer).
4. Let $s$ and $t$ be positive integers. Recall that $P_{s}$ is the path graph on $s$ vertices.
(a) For what values of $s$ and $t$ (if any) is $P_{s}$ an induced subgraph of $K_{t}$ ?
(b) For what values of $s$ and $t$ is $P_{s}$ a subgraph of $K_{t}$ ?
(c) For the values of $s$ and $t$ given in (b), provide a formula for the number of edges that need to be removed from $K_{t}$ to obtain $P_{s}$ as a subgraph.
5. Recall that an Eulerian walk in a graph is a walk that uses each edge exactly once. Recall that a connected graph has an Eulerian walk if and only if it has either 0 or 2 vertices of odd degree. What happens if, instead, we look for a walk that uses each edge exactly twice - once in each direction? Prove or disprove that every connected graph contains such a walk.
6. Prove Lemma 2.3: "Let $T$ be a tree with at least two vertices. Then $T$ has at least two leaves."
7. Prove Lemma 2.4: "Let $G$ be a connected graph. If $e$ is a bridge of $G$, then $G \backslash e$ has exactly two components."
8. A saturated hydrocarbon is a molecule $C_{m} H_{n}$ is which every carbon atom has four bonds, every hydrogen atom has one bond, and no sequence of bonds forms a cycle. Show that, for any positive integer $m$, the molecule $C_{m} H_{n}$ can exist if and only if $n=2 m+2$.
9. It is a famous theorem in graph theory that $K_{n}$ has $n^{n-2}$ spanning trees.
(a) Verify the theorem for $K_{2}, K_{3}$ and $K_{4}$.
(b) What about $K_{1}$ ?
(c) Can you think of strategies for proving the theorem in general?
