

Teaching/course surveys - fill in online

Recap: 4-color theorem (4CT)

- Kempe's erroneous proof
- Tait's unsuccessful approach

Theorem 6.15: Every loopless planar graph is 4-colorable if and only if every 3-connected cubic planar graph is 3-edge-colorable.

Let G be a plane graph. A proper k -face coloring is a function $\psi : F(G) \rightarrow S$ where $|S| = k$ and $\psi(f) \neq \psi(f')$ when f is adjacent to f' .

G is k -face-colorable if it has a proper k -face-coloring.

Theorem 6.17: A 3-connected cubic planar graph is 4-face-colorable if and only if it is 3-edge-colorable.

Proof of (\Rightarrow) Let G be a 3-connected cubic plane graph that has a proper 4-face-coloring $\psi : F(G) \rightarrow \mathbb{Z}_2 \times \mathbb{Z}_2$.

We define an edge-coloring ϕ .

For an edge e that is incident to faces f_1 and f_2 , let

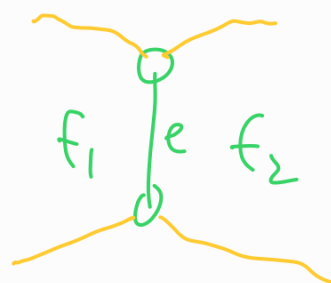
$$\phi(e) = \psi(f_1) + \psi(f_2).$$

$\mathbb{Z}_2 \times \mathbb{Z}_2 = \{ (0,0), (0,1), (1,0), (1,1) \}$
 forms group with addition
 (coordinate-wise mod 2)
 identity is $(0,0)$
 Let $a, b \in \mathbb{Z}_2 \times \mathbb{Z}_2$
 $a + b = (0,0)$ iff $a = b$ (*)

Then $\psi(f_1) \neq \psi(f_2)$, so $\varphi(e) \neq (0,0)$.

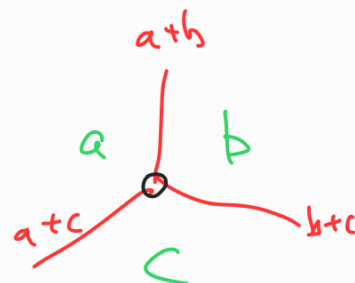
So $\varphi: E(G) \rightarrow (\mathbb{Z}_2 \times \mathbb{Z}_2) \setminus \{(0,0)\}$.

i.e. φ is a 3-edge-colouring.



It remains to show φ is proper.

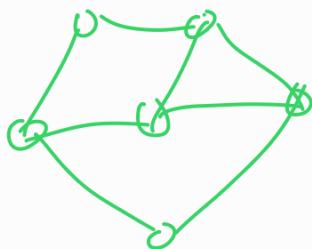
Since G is cubic, there are 3 faces incident with each vertex. Say the (distinct) colours these are assigned are a, b, c .



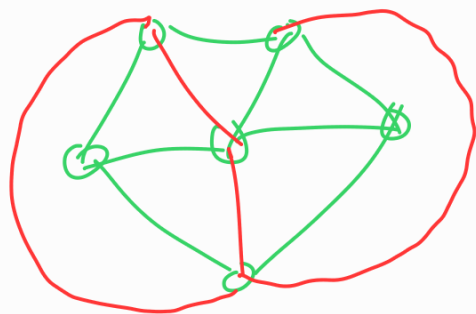
Since $\begin{cases} a \neq c \\ b \neq c \\ a \neq b \end{cases}$ we have $\begin{cases} a+b \neq c+b \\ b+a \neq c+a \\ a+c \neq b+c \end{cases}$, so $a+b, a+c$, and $b+c$ are distinct. \square

A plane graph is a triangulation if every face boundary, including the outer face boundary, is a cycle of length 3 (a triangle).

e.g.



A plane graph G



A triangulation H where G can be obtained from H by deleting the red edges.

Lemma 6.18 Every loopless planar graph is 4-colourable if and only if every triangulation is 4-colourable.

Let G be a connected plane graph.

Lemma 6.19: If G is a triangulation, then G^* is cubic.

Corollary 6.20: If G is cubic, then G^* is a triangulation.

Theorem 6.15: Every loopless planar graph is 4-colourable if and only if every 3-connected cubic planar graph is 3-edge-colourable.

Proof of Thm 6.15 (\Rightarrow)

Suppose every loopless planar graph is 4-colourable. Then every triangulation is 4-colourable by Lemma 6.18.

For every connected cubic plane graph G , G^* is a triangulation (Corollary 6.20).

Since the dual of every connected cubic plane graph G is 4-colourable, every connected cubic plane graph G is 4-face-colourable (by planar duality).

In particular, every 3-connected cubic plane graph is 4-face-colorable, so, by Thm 6.17, is also 3-edge-colorable, as required.

The reverse direction is similar. \square

Hadwiger's Conjecture

A graph G is k -chromatic if its chromatic number is at least k .

$$k\text{-chromatic: } \chi(G) \geq k$$

$$k\text{-colorable: } \chi(G) \leq k$$

Conjecture 6.23 (Hadwiger's Conjecture, 1943)

Let G be a loopless graph, and let k be a positive integer. If G is k -chromatic, then G has a K_k -minor.

The cases $k \leq 4$ are not too difficult.

$k=5$ was shown to be equivalent to 4CT

$k=6$ shown by Robertson, Seymour, Thomas (1993)

$k \geq 7$ open.